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## BLOCK-CLOSED SUBGRAPHS AND $q$ -PARTITIONS OF GRAPHS

*Abstract.* In this paper we introduce the concept of a block-closed subgraph of a graph and with help of this concept we continue the study of  $q$ -partitions of graphs which was initiated by Majcher and Płonka [2].

All graphs considered are finite undirected without loops and multiple edges.

Let  $G$  be a connected graph and let  $G_0$  be its induced subgraph. If  $G_0$  is connected and each block of  $G_0$  is a block of  $G$ , then  $G_0$  is called *block-closed*. Thus each block-closed subgraph of  $G$  is the union of some blocks of  $G$  which is connected. The empty graph  $K_0$  (with the empty vertex set) is also considered as a block-closed subgraph of any graph.

**THEOREM 1.** *All block-closed subgraphs of a connected graph  $G$  form a lattice  $LB(G)$  with respect to the ordering by inclusion.*

**Proof.** The intersection of two block-closed subgraphs  $G_1$  and  $G_2$  of  $G$  is evidently either empty or consists of one vertex (articulation point of  $G$ ), or is a block-closed subgraph of  $G$ . If it consists of one vertex, then we define the meet  $G_1 \wedge G_2$  as  $K_0$ . In other cases, the meet  $G_1 \wedge G_2$  is equal to the intersection of  $G_1$  and  $G_2$ . If the intersection of  $G_1$  and  $G_2$  is non-empty, then their union is connected, and thus it is a block-closed subgraph of  $G$ ; in this case, we define the join  $G_1 \vee G_2$  as this union. If the intersection of  $G_1$  and  $G_2$  is empty, then their union is disconnected. As  $G$  is connected, there exists a path connecting a vertex of  $G_1$  with a vertex of  $G_2$  and containing no edge of  $G_1$  and no edge of  $G_2$ . For each of such paths, the set of blocks of  $G$  containing at least one edge of that path is the same; otherwise, there would exist a circuit containing edges of different blocks. If we add all these blocks to the union of  $G_1$  and  $G_2$ , we obtain the least block-closed subgraph of  $G$  containing both  $G_1$  and  $G_2$  and this is  $G_1 \vee G_2$ . The greatest element of  $LB(G)$  is  $G$ , the least element is  $K_0$ .

The lattice  $LB(G)$  has properties analogous to those of the lattice of all subtrees of a tree [3].

**THEOREM 2.** *Each element of  $LB(G)$  is a join of atoms.*

**Proof.** Atoms of  $LB(G)$  are blocks of  $G$ . Each block-closed subgraph of  $G$  is the union of some blocks of  $G$ , which implies the assertion.

**THEOREM 3.** *If  $G$  contains at most one articulation point, then  $LB(G)$  is a Boolean algebra. If  $G$  contains more than one articulation point, then  $LB(G)$  is not modular.*

**Proof.** If  $G$  is an empty graph, then  $LB(G)$  consists of one element. If  $G$  is non-empty without articulation points, then  $LB(G)$  consists of two elements  $K_0$  and  $G$  and is a Boolean algebra with one generator. If  $G$  contains exactly one articulation point, then the union of any blocks of  $G$  is a connected graph, and thus a block-closed subgraph of  $G$ . As all block-closed subgraphs of  $G$  are unions of blocks of  $G$ , there is a one-to-one correspondence between subsets of the set of blocks of  $G$  and block-closed subgraphs of  $G$  such that each block-closed subgraph of  $G$  is the union of all blocks of the corresponding set. This correspondence is evidently a lattice isomorphism of the Boolean algebra of all subsets of the set of blocks of  $G$  onto  $LB(G)$ , and hence  $LB(G)$  is a Boolean algebra.

Now suppose that  $G$  has at least two articulation points. Then there exist blocks  $B_0, B_1, B_2$  of  $G$  such that  $B_0$  and  $B_1$  have a common vertex  $a_1$ , the blocks  $B_0$  and  $B_2$  have a common vertex  $a_2$ , and  $B_1$  and  $B_2$  have no common vertex. Evidently,  $a_1$  and  $a_2$  are articulation points of  $G$  and  $a_1 \neq a_2$ . Let  $H_0$  be the union of  $B_0$  and  $B_1$ , and let  $H$  be the union of  $B_0, B_1, B_2$ . The graphs  $B_0, B_1, B_2, H_0, H$  belong to  $LB(G)$ . We have

$$B_1 \wedge B_2 = H_0 \wedge B_2 = K_0, \quad B_1 \vee B_2 = H_0 \vee B_2 = H,$$

and thus the elements  $K_0, B_1, B_2, H_0, H$  form a "forbidden pentagon" by the definition of modularity, and hence  $LB(G)$  is not modular.

In the sequel we shall study trees. A graph  $K_1$  consisting of one vertex and the empty graph  $K_0$  are also considered as trees.

The tree of blocks [1] of a connected graph  $G$  is the tree whose vertex set is the union of the set of blocks of  $G$  and the set of articulation points of  $G$  and in which each vertex being a block of  $G$  is adjacent to all vertices being articulation points of  $G$  contained in this block.

Let  $\mathcal{T}_0$  be the class consisting of all finite trees  $T$  with the property that the distance between any two terminal vertices of  $T$  is even. The graphs  $K_0$  and  $K_1$  also belong to  $\mathcal{T}_0$ .

**THEOREM 4.** *Let  $T$  be a finite tree. Then the following two assertions are equivalent:*

- (i)  $T$  is isomorphic to the tree of blocks of some graph  $G$ .
- (ii)  $T \in \mathcal{T}_0$ .

**Proof.** (i)  $\Rightarrow$  (ii). Let  $T$  be isomorphic to the tree of blocks of some graph  $G$ ; we may identify it with this tree. Then the vertex set of  $T$  is the

union of two disjoint sets  $V_0$  and  $V_1$ , where  $V_0$  is the set of blocks of  $G$ , and  $V_1$  is the set of articulation points of  $G$ . Two vertices can be adjacent only if one of them belongs to  $V_0$  and the other to  $V_1$ . Hence the distance between any two vertices of  $V_0$  is even. The terminal vertices of  $T$  are always blocks of  $G$ , and therefore they belong to  $V_0$  and the distance between any two of them is even, i.e.,  $T \in \mathcal{T}_0$ .

(ii)  $\Rightarrow$  (i). Let  $T \in \mathcal{T}_0$ . Choose an arbitrary terminal vertex  $u$  of  $T$ . By  $V_0$  (or  $V_1$ ) denote the set of all vertices of  $T$  whose distance from  $u$  is even (or odd, respectively). As every tree is a bipartite graph, we have  $V_0 \cap V_1 = \emptyset$ . As  $T \in \mathcal{T}_0$ , these sets are independent of the choice of  $u$ . Now we add new edges to  $T$ . Two vertices of  $T$  are joined by a new edge if and only if they both belong to  $V_1$  and their distance in  $T$  is 2. The graph thus obtained is denoted by  $G$ . Each block of  $G$  is a clique whose vertex set consists of one vertex of  $V_0$  and all vertices of  $V_1$  which are adjacent to it in  $T$ . Evidently, the tree of blocks of  $G$  is isomorphic to  $T$ .

Now let  $T \in \mathcal{T}_0$ . We define a  $\mathcal{T}_0$ -subtree of  $T$  as a subtree of  $T$ , all of whose terminal vertices have even distances from terminal vertices of  $T$ . Evidently, each  $\mathcal{T}_0$ -subtree of  $T$  belongs to  $\mathcal{T}_0$ . We admit also trees having only one vertex and the empty graph  $K_0$ .

**THEOREM 5.** *All  $\mathcal{T}_0$ -subtrees of a tree  $T \in \mathcal{T}_0$  form a lattice  $LT_0(T)$  with respect to ordering by inclusion.*

**Proof.** For any two  $\mathcal{T}_0$ -subtrees  $T_1$  and  $T_2$  of  $T$  we define the meet  $T_1 \wedge T_2$  and the join  $T_1 \vee T_2$ . The intersection of  $T_1$  and  $T_2$  is either empty or is a subtree of  $T$ . If it is a subtree of  $T$  not belonging to  $LT_0(T)$ , then  $T_1 \wedge T_2$  is obtained from it by deleting all terminal vertices whose distance from terminal vertices of  $T$  is odd; such a tree is evidently the greatest  $\mathcal{T}_0$ -subtree of  $T$  contained in both  $T_1$  and  $T_2$ . In other cases,  $T_1 \wedge T_2$  is the intersection of  $T_1$  and  $T_2$ . If the intersection of  $T_1$  and  $T_2$  is non-empty, then the union of  $T_1$  and  $T_2$  is a tree; each terminal vertex of this tree is a terminal vertex of  $T_1$  or of  $T_2$ , and thus this union is in  $LT_0(T)$  and  $T_1 \vee T_2$  is equal to it. If the intersection of  $T_1$  and  $T_2$  is empty, then there exists exactly one path in  $T$  connecting a vertex of  $T_1$  with a vertex of  $T_2$  and containing no edge of  $T_1$  and no edge of  $T_2$ . By adding it to the union of  $T_1$  and  $T_2$  we obtain a tree with the property that each terminal vertex of it is a terminal vertex of  $T_1$  or of  $T_2$ . Therefore, this tree belongs to  $LT_0(T)$  and it is the smallest tree of  $LT_0(T)$  which contains both  $T_1$  and  $T_2$ ; hence it is  $T_1 \vee T_2$ . The greatest element of  $LT_0(T)$  is  $T$ , the least is  $K_0$ .

**THEOREM 6.** *Let  $G$  be a graph, and  $T$  be its tree of blocks. Then*

$$LB(G) \cong LT_0(T).$$

**Proof.** This is an easy consequence of Theorems 4 and 5.

COROLLARY 1. *If the diameter of  $T$  is at most 2, then  $LT_0(T)$  is a Boolean algebra. If it is greater than 2, then  $LT_0(T)$  is not modular.*

THEOREM 7. *Let the lattice  $LB(G)$  of a graph  $G$  be given as an abstract lattice. Then the tree of blocks of  $G$  can be reconstructed uniquely up to isomorphism.*

Proof. Let  $LB(G)$  be given. Let  $B_1, \dots, B_n$  be its atoms; they are blocks of  $G$ . If  $i, j, k$  are pairwise different numbers from the set  $\{1, \dots, n\}$  and  $(B_i \vee B_j) \wedge B_k = B_k$ , then each block-closed subgraph of  $G$  containing  $B_i$  and  $B_j$  contains also  $B_k$ , and therefore  $B_i$  and  $B_j$  have no common vertex. Hence  $B_i$  and  $B_j$  have a common vertex if and only if

$$(B_i \vee B_j) \wedge B_k = \emptyset \quad \text{for all } k \in \{1, \dots, n\} - \{i, j\}.$$

Now we can construct the family  $\mathcal{S}$  of maximal sets  $S$  of blocks of  $G$  with the property that any two blocks from  $S$  have a common vertex. From the properties of blocks it follows that for each  $S \in \mathcal{S}$  there exists an articulation point of  $G$  which is common to all blocks of  $S$ . Conversely, to each articulation point of  $G$  a set from  $\mathcal{S}$  can be assigned which consists of blocks containing this point. This correspondence is evidently one-to-one, therefore we may reconstruct the articulation points of  $G$  and determine, for each block, which articulation points it contains; this means that we are able to reconstruct the tree of blocks of  $G$ .

COROLLARY 2. *Let the lattice  $LT_0(T)$  of a tree  $T \in \mathcal{T}_0$  be given as an abstract lattice. Then  $T$  can be reconstructed uniquely up to isomorphism.*

Now we turn our attention back to  $q$ -partitions.

In [2], an induced subgraph  $Q$  of a connected graph  $G$  is called a *quasi-component* of  $G$  if the number of vertices of  $Q$  which are adjacent in  $G$  to a vertex not belonging to  $Q$  is at most one. If there is one such vertex, it is called the *control point* of  $Q$  and denoted by  $c(Q)$ . It is proved in [2] that every quasi-component of a graph  $G$  is either a one-vertex subgraph of  $G$  or a connected union of blocks of  $G$ . A partition  $P$  of the vertex set  $V(G)$  of  $G$  is called a  *$q$ -partition* of  $G$  if each class of  $P$  induces a quasi-component of  $G$ .

It has been proved that all  $q$ -partitions of a graph  $G$  form a lattice. The ordering in this lattice is defined as follows. The symbol  $[a]_P$  denotes the class of  $P$  which contains vertex  $a$ . We have  $P_1 \leq P_2$  if and only if  $[a]_{P_1} \subseteq [a]_{P_2}$  for each vertex  $a$  of  $G$ .

In what follows we study the interconnections between block-closed subgraphs and  $q$ -partitions of  $G$ .

Let  $G$  be a graph and let  $H$  be its block-closed subgraph. Let  $P(H)$  be the partition of  $V(G)$  such that two vertices of  $G$  belong to the same class of  $P(H)$  if and only if they are connected by a path in  $G$  which does not contain any edge of  $H$ .

**THEOREM 8.** *For each block-closed subgraph  $H$  of  $G$  the partition  $P(H)$  is a  $q$ -partition of  $G$ .*

**Proof.** The subgraphs of  $G$  induced by the classes of  $P(H)$  are connected components of the graph  $G'$  obtained from  $G$  by deleting all edges of  $H$ . Hence each of them is either a union of some blocks of  $G$  not belonging to  $H$  or a one-vertex graph whose vertex belongs to  $H$ . In the second case it is evidently a quasi-component of  $G$ . Let the first case occur and let  $Q$  be such a subgraph. Suppose that there exist two vertices  $a$  and  $b$  of  $Q$  which are incident with edges of  $G$  not belonging to  $Q$ , i.e., edges of  $H$ . Since each quasi-component of  $G$  is connected (see [2]), there exists a path  $P_1$  in  $Q$  connecting  $a$  and  $b$ . As  $a$  and  $b$  are incident with edges of  $H$ , they belong to  $H$ . As  $H$  is connected, there exists a path  $P_2$  in  $H$  connecting  $a$  and  $b$ . The union of  $P_1$  and  $P_2$  is a circuit containing edges of  $H$  and edges of  $Q$ . Thus the edges of  $P_1$  belong to the same block of  $G$  as those of  $P_2$ . This block contains edges of  $H$ , but not all of them; hence the blocks of  $H$  containing these edges are not blocks of  $G$  and  $H$  is not block-closed, which is a contradiction. Hence  $Q$  is a quasi-component of  $G$  and  $P(H)$  is a  $q$ -partition of  $G$ .

**THEOREM 9.** *Let  $P$  be a  $q$ -partition of  $G$ . Then there exists a block-closed subgraph  $H$  of  $G$  such that  $P = P(H)$ .*

**Proof.** Let  $H$  be the subgraph of  $G$  whose edge set is the set of all edges joining pairs of vertices from different classes of  $P$  and whose vertex set is the set of end vertices of these edges. Suppose that  $H$  is not connected. Then there exist vertices  $a$  and  $b$  in  $H$  which are connected by a path  $R$  in  $G$ , none of whose edges is in  $H$ . As each edge joining vertices of different classes of  $P$  lies in  $H$ , all edges of  $R$  lie in the subgraph of  $G$  induced by one class  $Q$  of  $P$ . But then  $a$  and  $b$  are two vertices of  $Q$  incident with edges not belonging to  $Q$  and  $Q$  is not a quasi-component of  $G$ , which is a contradiction. Hence  $H$  is connected. Suppose that there exists a block  $B$  of  $H$  which is not a block of  $G$ . Then  $B$  is a proper subgraph of a block  $B'$  of  $G$ . Let  $e_1$  be an edge of  $B$ , and  $e_2$  be an edge of  $B'$  not belonging to  $B$ . As both  $e_1$  and  $e_2$  belong to  $B'$ , there exists a circuit  $C$  in  $B'$  containing both  $e_1$  and  $e_2$ . The circuit  $C$  contains a path (containing  $e_2$ ) which connects two vertices  $a$  and  $b$  of  $B$  and whose edges are not contained in  $H$ . Then we obtain the same contradiction as in the preceding case. Hence  $H$  is a block-closed subgraph of  $G$ . Evidently,  $P = P(H)$ .

**THEOREM 10.** *The lattice of  $q$ -partitions of a graph  $G$  is dually isomorphic to  $LB(G)$ .*

**Proof.** From Theorems 8 and 9 it is clear that there exists a one-to-one correspondence between the graphs  $H$  from  $LB(G)$  and the  $q$ -partitions  $P(H)$  of  $G$ . If  $H_1$  and  $H_2$  are in  $LB(G)$  and  $H_1$  is a subgraph of  $H_2$ , then each edge

not belonging to  $H_2$  does not belong to  $H_1$  and thus  $[a]_{P(H_2)} \subseteq [a]_{P(H_1)}$  for each vertex  $a$  of  $G$ . Hence  $P(H_2) \leq P(H_1)$ . This implies the assertion.

**COROLLARY 3.** *Each element of the lattice of all  $q$ -partitions of  $G$  is a meet of dual atoms.*

**COROLLARY 4.** *If  $G$  contains at most one articulation point, then the lattice of all  $q$ -partitions of  $G$  is a Boolean algebra; otherwise, it is not modular.*

**COROLLARY 5.** *Let the lattice of all  $q$ -partitions of a graph  $G$  be given as an abstract lattice. Then the tree of blocks of  $G$  can be reconstructed uniquely up to isomorphism.*

### References

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