# A SHORT NOTE ON $L_{C B A}$-FUZZY LOGIC WITH A NON-ASSOCIATIVE CONJUNCTION 

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#### Abstract

We significantly simplify the axiomatic system $L_{C B A}$ for fuzzy logic with a non-associative conjunction.


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Several investigations in probability theory and the theory of expert systems show that it is important to search for some reasonable generalizations of fuzzy logics having a non-associative conjunction, see $[5,6,8,9,10]$.

In the paper [1] Botur and Halaš introduced and described a non-associative fuzzy logic $L_{C B A}$ having as an equivalent algebraic semantics lattices with section antitone involutions satisfying the contraposition law, so-called commutative basic algebras. The variety of commutative basic algebras was intensively studied in several recent papers and includes the class of MV-algebras. For more details see the book [2].

In [1] Botur and Halaš introduced a non-associative fuzzy logic $L_{C B A}$ by the following nine axioms:
(i) $x \rightarrow(y \rightarrow x)=1$
(ii) $((x \rightarrow y) \rightarrow y) \rightarrow((y \rightarrow x) \rightarrow x)=1$
(iii) $(\neg x \rightarrow \neg y) \rightarrow(y \rightarrow x)=1$
(iv) $\neg \neg x \rightarrow x=1$
(v) $(x \rightarrow y) \rightarrow(x \rightarrow \neg \neg y)=1$
(vi) $x \rightarrow x=1$
(vii) $(x=1 \& x \rightarrow y=1) \Rightarrow y=1$
(viii) $x \rightarrow y=1 \Rightarrow(y \rightarrow z) \rightarrow(x \rightarrow z)=1$
(ix) $(x \rightarrow y=1 \& y \rightarrow x=1) \Rightarrow x=y$.

In the main theorem we significantly simplify the axiomatic system $L_{C B A}$ for fuzzy logic with a non-associative conjunction. We show that three axioms can be omitted.

Theorem 1. Identities (iv), (v), (vi) follows from identities (i), (ii), (iii) and from quasiidentities (vii), (viii), (ix).

Proof. First, we show that (vi) is redundant. Using (viii) on (i) we easily obtain

$$
\begin{equation*}
((y \rightarrow x) \rightarrow x) \rightarrow(x \rightarrow x)=1 \tag{1}
\end{equation*}
$$

Using (viii) on (ii) we obtain

$$
\begin{equation*}
(((y \rightarrow x) \rightarrow x) \rightarrow(x \rightarrow x)) \rightarrow(((x \rightarrow y) \rightarrow y) \rightarrow(x \rightarrow x))=1 \tag{2}
\end{equation*}
$$

Applying (vii) on (1) and (2) we get

$$
\begin{equation*}
((x \rightarrow y) \rightarrow y)) \rightarrow(x \rightarrow x)=1 \tag{3}
\end{equation*}
$$

Further, using (viii) on (i) in the form $y \rightarrow((x \rightarrow y) \rightarrow y)=1$ we derive

$$
\begin{equation*}
(((x \rightarrow y) \rightarrow y) \rightarrow(x \rightarrow x)) \rightarrow(y \rightarrow(x \rightarrow x))=1 \tag{4}
\end{equation*}
$$

Applying (vii) on (3) and (4) we get $y \rightarrow(x \rightarrow x)=1$. From the last identity, where $y:=1$, using (vii), we obtain (vi).

Now, we show that (iv) is redundant. From (i), using (vi), we easily derive that

$$
\begin{equation*}
x \rightarrow 1=1 \tag{5}
\end{equation*}
$$

Putting $y:=1$ in (ii) and applying (5) we get

$$
1 \rightarrow((1 \rightarrow x) \rightarrow x)=1
$$

which, by (vii), give us

$$
\begin{equation*}
(1 \rightarrow x) \rightarrow x=1 \tag{6}
\end{equation*}
$$

From this and from (i) in the form $x \rightarrow(1 \rightarrow x)=1$, we obtain

$$
\begin{equation*}
1 \rightarrow x=x \tag{7}
\end{equation*}
$$

by (ix). From (iii) and (viii) we infer

$$
\begin{equation*}
((y \rightarrow x) \rightarrow z) \rightarrow((\neg x \rightarrow \neg y) \rightarrow z)=1 \tag{8}
\end{equation*}
$$

from (i) and (viii) we obtain

$$
\begin{equation*}
((y \rightarrow x) \rightarrow z) \rightarrow(x \rightarrow z)=1, \tag{9}
\end{equation*}
$$

and from (ii) and (ix) we have

$$
\begin{equation*}
(x \rightarrow y) \rightarrow y=(y \rightarrow x) \rightarrow x . \tag{10}
\end{equation*}
$$

Putting $y:=1$ in (iii) and applying (7) we get

$$
\begin{equation*}
(\neg x \rightarrow \neg 1) \rightarrow x=1 \tag{11}
\end{equation*}
$$

From this, for $x:=\neg 1$, we get immediately $(\neg \neg 1 \rightarrow \neg 1) \rightarrow \neg 1=1$. From (i), where $x:=\neg 1$ and $y:=\neg \neg 1$, we obtain $\neg 1 \rightarrow(\neg \neg 1 \rightarrow \neg 1)=1$. Applying (ix) to the last two equations we conclude

$$
\begin{equation*}
\neg \neg 1 \rightarrow \neg 1=\neg 1 . \tag{12}
\end{equation*}
$$

Applying (vii) to (11) and to (9), where $y:=\neg x, x:=\neg 1$ and $z:=x$, we derive

$$
\begin{equation*}
\neg 1 \rightarrow x=1 \tag{13}
\end{equation*}
$$

Now, put $\neg 1$ instead of $x$ and $x \rightarrow \neg \neg 1$ instead of $y$ in (10) to obtain

$$
(\neg 1 \rightarrow(x \rightarrow \neg \neg 1)) \rightarrow(x \rightarrow \neg \neg 1)=((x \rightarrow \neg \neg 1) \rightarrow \neg 1) \rightarrow \neg 1 .
$$

The left hand side of the last identity can be reduced to $x \rightarrow \neg \neg 1$, using (13) and (7). The right hand side of the last identity is equal to 1 , by (9), where $y:=x$, $x:=\neg \neg 1$ and $z:=\neg 1$, using (12). Therefore, we have

$$
\begin{equation*}
x \rightarrow \neg \neg 1=1 . \tag{14}
\end{equation*}
$$

From (14) we derive easily that $1 \rightarrow \neg \neg 1=1$ and from (5) we derive easily that $\neg \neg 1 \rightarrow 1=1$, whence, by (ix), we conclude

$$
\begin{equation*}
\neg \neg 1=1 \text {. } \tag{15}
\end{equation*}
$$

Putting $y:=\neg x, x:=\neg 1$ and $z:=x$ in (8) and applying (11) and (vii) we get

$$
(\neg \neg 1 \rightarrow \neg \neg x) \rightarrow x=1
$$

which, by (15) and (5), give us (iv).
Finally, we show that also (v) is redundant. Using (vii) on (iii) in the form $(\neg \neg \neg x \rightarrow \neg x) \rightarrow(x \rightarrow \neg \neg x)=1$ and on (iv) in the form $\neg \neg \neg x \rightarrow x=1$ we get $x \rightarrow \neg \neg x=1$. This with (iv) give us by (ix)

$$
x=\neg \neg x .
$$

The last identity together with (vi) in the form $(x \rightarrow y) \rightarrow(x \rightarrow y)=1$ proves (v).

Among the recent results on simplification of axiomatic systems related to fuzzy logics belong works of Cintula (based on a proof, see [3]) and Lehmke (applying a solver he programmed, see [7]). Both of them have shown the redundancy of one of the axioms of Hájek's axiomatization of BL-logics [4].

## Conclusions

We briefly recalled what a non-associative fuzzy logic $L_{C B A}$ is and where it can be useful. Then we significantly simplify the axiomatic system $L_{C B A}$ for fuzzy logic with a non-associative conjunction as given in [1]. We show that three axioms can be omitted. Note that generally, the removal of the redundant axioms from a axiomatic system is desirable because it simplifies definitions and shortens proofs.

Note that although the independence of axioms is a highly desirable property, superfluous axioms still could be of use (as tautologies in the considered theory) when proving results in the considered theory, making proofs more transparent.

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